CHAPTER 2 SECTION 2.4

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Recursive Definitions

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- A recursive definition is one in which the item being defined is included as part of the definition
 Also called an inductive definition
- 2 parts to a recursive definition
 - 1. A basis, where some simple cases (1 or more) of the item being defined are explicitly given
 - 2. A recursive or inductive step, where new cases of the item being defined are given in terms of previous cases

Recursively Defined Sequences

Sequence S

- A list of objects that are enumerated in some order
- S(k) denotes the k^{th} object in the sequence
- Defining a sequence recursively
 - First, name one or more base cases
 - Then, define later values in terms of earlier ones
- Example 29: Sequence S is defined recursively by

1.
$$S(1) = 2$$

2. $S(n) = 2S(n-1)$ for $n \ge 2$

Recursively Defined Sequences

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Fibonacci Sequence of Numbers

- F(1) = 1 F(2) = 1F(n) = F(n-2) + F(n-1) for n > 2
- What is the basis?
- Write the first 10 values of the sequence.

Fibonacci Numbers

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- □ Problem 14: Prove the given property of the Fibonacci numbers directly from the definition. F(n+3) = 2F(n+1) + F(n) for $n \ge 1$
- □ Recall that the definition (recursive part) is F(k) = F(k-2) + F(k-1)
- □ Since we're trying to show a property F(n+3), substitute n+3 into the definition for k (k=n+3). F(n+3) = F(n+3-2) + F(n+3-1) = F(n+1) + F(n+2)

Fibonacci Numbers

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- Now, prove the same formula using the 2nd principle of induction (Problem 20). P(n): F(n+3) = 2F(n+1) + F(n) for $n \ge 1$
- Step 1: Base Case
 - Since we are using 2 previous values to compute the next value, use 2 base cases

Problem 20 continued

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Step 2.a: Assume P(r): F(r+3) = 2F(r+1) + F(r) for $1 \le r \le k$ What does this mean? We can assume $F(1+3) = 2F(1+1) + F(1) \implies F(4) = 2F(2) + F(1)$ $F(2+3) = 2F(2+1) + F(2) \Rightarrow F(5) = 2F(3) + F(2)$... $\blacksquare F(k-1+3) = 2F(k-1+1) + F(k-1) \Longrightarrow$ F(k+2) = 2F(k) + F(k-1)F(k+3) = 2F(k+1) + F(k)

Problem 20 continued

Step 2.b: Prove

P(k+1): F(k+1+3) ?= 2F(k+1+1) + F(k+1) $\Rightarrow F(k+4) ?= 2F(k+2) + F(k+1)$

• Start with the definition of Fibonacci numbers E(x) = E(x) + E(x-1)

F(n) = F(n-2) + F(n-1)

Since we are trying to prove F(k+4), substitute k+4 for n in the definition formula

F(k+4) = F(k+4-2) + F(k+4-1) = F(k+2) + F(k+3)

Now, do we have any information we can use about F(k+2) and F(k+3)?

Fibonacci Sequence

Example 31: Prove that in the Fibonacci sequence

F(n+4) = 3F(n+2) - F(n) for all $n \ge 1$

Proof by induction

Fibonacci Sequence

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Prove the formula without induction

F(n + 4) = 3F(n + 2) - F(n) for all $n \ge 1$

Use the recurrence relation from the definition of Fibonacci numbers

F(n) = F(n - 2) + F(n - 1)

Recursively Defined Sets

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- A sequence is a collection of objects which have a specific order
- A set is a collection of object with no order imposed
- Example: A recursive definition for the set of propositional wffs
 - 1. Any statement letter is a wff.
 - 2. If P and Q are wffs, so are (P v Q), (P \land Q), (P \rightarrow Q), (P'), and (P \leftrightarrow Q)
- □ Show how to build $((A \lor (B')) \rightarrow C)$

Recursively Defined Sets

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- Example 34: The set of all (finite-length) strings of symbols over a finite alphabet A is denoted by A*. The recursive definition of A* is

1. The empty string λ (the string with no symbols) belongs to A^* .

2. Any single member of A belongs to A^* .

3. If x and y are strings in A^* , so is xy, the concatenation of strings x and y.

□ If x = 1011 and y = 001, write the strings xy, yx, and $yx\lambda x$.

Recursively Defined Set

Practice 17

Give a recursive definition for the set of all *binary* strings that are **palindromes**.

A palindrome is a string that reads the same forwards and backwards.

Backus-Naur Form

- BNF notation allows you to recursively define a set of strings
 - Angle brackets < > indicate items that are defined in terms of other items
 - Items without brackets cannot be further broken down
 - The vertical line | means or

Backus-Naur Form

A BNF definition of an identifier
<identifier> ::= <letter> | <identifier><letter> | <identifier><letter> | <identifier><digit>
<letter> ::= a | b | c | ... | z
<digit> ::= 1 | 2 | ... | 9
How would the identifier tmn1 be built from

How would the identifier *tmp*1 be built from the definition?

Recursively Defined Operations

- Some operations can also be defined recursively
 - Example 36: Exponentiation operation aⁿ on a nonzero real number a, where n is a nonnegative integer
 - Recursive definition

1.
$$a^0 = 1$$

2. $a^n = (a^{n-1})a$ for $n \ge 1$

Recursively Defined Operations

Practice 18

- Let x be a string over some alphabet
- Give a recursive definition for the operation x^n (concatenation of x with itself n times) for $n \ge 1$.

Recursively Defined Algorithms

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- Write a computer algorithm to calculate S(n) from Example 29
 - 1. S(1) = 2
 - 2. S(n) = 2S(n-1) for $n \ge 2$
- Iterative vs. recursive

Binary Search

Input

- List of items sorted in nondecreasing order
- An item x which you would like to find

Basic Idea

- Compare x to the middle item in list
- If it matches, you' re done.
- If x is less than middle item
 - Search first half of list
- If x is greater than middle item
 - Search second half of list

Binary Search Algorithm

BinarySearch(list L; integer i; integer j; itemtype x) // searches sorted list L from L[i] to L[i] for item x **if** (i > j) **then** write ("not found") else find the index k of the middle item in the list L[i]-L[j] if x = middle item then write("found") else if x < middle item then BinarySearch(L, I, k-1, x) else BinarySearch(L, k+1, j, x) end if end if end if end function BinarySearch

Binary Search

Apply the binary search algorithm to the list
3, 7, 8, 10, 14, 18, 22, 34