

# CHAPTER 2

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## Section 2.4

# Recursive Definitions

- A **recursive definition** is one in which the item being defined is included as part of the definition
  - Also called an **inductive definition**
- 2 parts to a recursive definition
  1. A basis, where some simple cases (1 or more) of the item being defined are explicitly given
  2. A recursive or inductive step, where new cases of the item being defined are given in terms of previous cases

# Recursively Defined Sequences

- Sequence S
  - A list of objects that are enumerated in some order
  - $S(k)$  denotes the  $k^{\text{th}}$  object in the sequence
- Defining a sequence recursively
  - First, name one or more base cases
  - Then, define later values in terms of earlier ones
- Example 29: Sequence S is defined recursively by
  1.  $S(1) = 2$
  2.  $S(n) = 2S(n-1)$  for  $n \geq 2$

# Recursively Defined Sequences

- Fibonacci Sequence of Numbers

$$F(1) = 1$$

$$F(2) = 1$$

$$F(n) = F(n-2) + F(n-1) \quad \text{for } n > 2$$

- What is the basis?
- Write the first 10 values of the sequence.

# Fibonacci Numbers

- Problem 14: Prove the given property of the Fibonacci numbers directly from the definition.

$$F(n+3) = 2F(n+1) + F(n) \text{ for } n \geq 1$$

- Recall that the definition (recursive part) is

$$F(k) = F(k-2) + F(k-1)$$

- Since we're trying to show a property  $F(n+3)$ , substitute  $n+3$  into the definition for  $k$  ( $k=n+3$ ).

$$F(n+3) = F(n+3-2) + F(n+3-1) = F(n+1) + F(n+2)$$

# Fibonacci Numbers

- Now, prove the same formula using the 2<sup>nd</sup> principle of induction (Problem 20).

$$P(n): F(n + 3) = 2F(n + 1) + F(n) \text{ for } n \geq 1$$

- Step 1: Base Case
  - Since we are using 2 previous values to compute the next value, use 2 base cases

# Problem 20 continued

- Step 2.a: Assume

$$P(r): F(r+3) = 2F(r+1) + F(r) \text{ for } 1 \leq r \leq k$$

- What does this mean? We can assume

- $F(1+3) = 2F(1+1) + F(1) \Rightarrow F(4) = 2F(2) + F(1)$
- $F(2+3) = 2F(2+1) + F(2) \Rightarrow F(5) = 2F(3) + F(2)$
- ...
- $F(k-1+3) = 2F(k-1+1) + F(k-1) \Rightarrow$   

$$F(k+2) = 2F(k) + F(k-1)$$
- $F(k+3) = 2F(k+1) + F(k)$

# Problem 20 continued

- Step 2.b: Prove

$$P(k+1): F(k+1+3) \stackrel{?}{=} 2F(k+1+1) + F(k+1)$$

$$\Rightarrow F(k+4) \stackrel{?}{=} 2F(k+2) + F(k+1)$$

- Start with the definition of Fibonacci numbers

$$F(n) = F(n-2) + F(n-1)$$

- Since we are trying to prove  $F(k+4)$ , substitute  $k+4$  for  $n$  in the definition formula

$$F(k+4) = F(k+4-2) + F(k+4-1) = F(k+2) + F(k+3)$$

- Now, do we have any information we can use about  $F(k+2)$  and  $F(k+3)$ ?



# Fibonacci Sequence

- Example 31: Prove that in the Fibonacci sequence

$$F(n + 4) = 3F(n + 2) - F(n) \quad \text{for all } n \geq 1$$

- Proof by induction

# Fibonacci Sequence

- Prove the formula without induction

$$F(n + 4) = 3F(n + 2) - F(n) \quad \text{for all } n \geq 1$$

- Use the recurrence relation from the definition of Fibonacci numbers

$$F(n) = F(n - 2) + F(n - 1)$$

# Recursively Defined Sets

- A **sequence** is a collection of objects which have a specific order
- A **set** is a collection of object with no order imposed
- Example: A recursive definition for the set of propositional wffs
  1. Any statement letter is a wff.
  2. If  $P$  and  $Q$  are wffs, so are  $(P \vee Q)$ ,  $(P \wedge Q)$ ,  $(P \rightarrow Q)$ ,  $(P')$ , and  $(P \leftrightarrow Q)$
- Show how to build  $((A \vee (B')) \rightarrow C)$

# Recursively Defined Sets

- Example 34: The set of all (finite-length) strings of symbols over a finite alphabet  $A$  is denoted by  $A^*$ . The recursive definition of  $A^*$  is
  1. The empty string  $\lambda$  (the string with no symbols) belongs to  $A^*$ .
  2. Any single member of  $A$  belongs to  $A^*$ .
  3. If  $x$  and  $y$  are strings in  $A^*$ , so is  $xy$ , the concatenation of strings  $x$  and  $y$ .
- If  $x = 1011$  and  $y = 001$ , write the strings  $xy$ ,  $yx$ , and  $yx\lambda x$ .

# Recursively Defined Set

- Practice 17
  - Give a recursive definition for the set of all *binary* strings that are **palindromes**.
    - A **palindrome** is a string that reads the same forwards and backwards.

# Backus-Naur Form

- BNF notation allows you to recursively define a set of strings
  - Angle brackets  $\langle \rangle$  indicate items that are defined in terms of other items
  - Items without brackets cannot be further broken down
  - The vertical line  $|$  means or

# Backus-Naur Form

- A BNF definition of an identifier
  - $\langle \text{identifier} \rangle ::= \langle \text{letter} \rangle \mid \langle \text{identifier} \rangle \langle \text{letter} \rangle \mid \langle \text{identifier} \rangle \langle \text{digit} \rangle$
  - $\langle \text{letter} \rangle ::= a \mid b \mid c \mid \dots \mid z$
  - $\langle \text{digit} \rangle ::= 1 \mid 2 \mid \dots \mid 9$
- How would the identifier *tmp1* be built from the definition?

# Recursively Defined Operations

- Some operations can also be defined recursively
  - Example 36: Exponentiation operation  $a^n$  on a nonzero real number  $a$ , where  $n$  is a nonnegative integer
  - Recursive definition
    1.  $a^0 = 1$
    2.  $a^n = (a^{n-1})a$  for  $n \geq 1$



# Recursively Defined Operations

- Practice 18
  - Let  $x$  be a string over some alphabet
  - Give a recursive definition for the operation  $x^n$  (concatenation of  $x$  with itself  $n$  times) for  $n \geq 1$ .

# Recursively Defined Algorithms

- Write a computer algorithm to calculate  $S(n)$  from Example 29
  1.  $S(1) = 2$
  2.  $S(n) = 2S(n-1)$  for  $n \geq 2$
- Iterative vs. recursive

# Binary Search

- Input
  - List of items sorted in nondecreasing order
  - An item  $x$  which you would like to find
- Basic Idea
  - Compare  $x$  to the middle item in list
  - If it matches, you're done.
  - If  $x$  is less than middle item
    - Search first half of list
  - If  $x$  is greater than middle item
    - Search second half of list

# Binary Search Algorithm

```
BinarySearch(list L; integer i; integer j; itemtype x)
// searches sorted list L from L[i] to L[j] for item x
  if (i > j) then
    write ("not found")
  else
    find the index k of the middle item in the list L[i]-L[j]
    if x = middle item then
      write("found")
    else
      if x < middle item then
        BinarySearch(L, i, k-1, x)
      else
        BinarySearch(L, k+1, j, x)
      end if
    end if
  end if
end function BinarySearch
```

# Binary Search

- Apply the binary search algorithm to the list  
3, 7, 8, 10, 14, 18, 22, 34