

# Network Metrics

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# Network Metrics

- Types of Errors
- Performance Measures

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## Reliability

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- Most users expect the network to deliver their data to the Application Layer without error
  - All the data that was sent is received (“all received”)
  - Within a unit of data, order is preserved (“in order”)
  - The data is not garbled (“no bit errors”)
- This means that lower layers must, as far as is feasible, try to prevent errors, and detect and correct errors that do occur

## What kinds of errors might occur?

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- Bit errors
  - random or bursts
  - errors in data can often be fixed
  - errors in header info can lead to strange effects
- Packet-level errors
  - delayed due to congestion
  - lost/mis-routed
  - out of order
- Link and node failures
  - short and long-term
  - unexpected start-up problems

Detection and correction?  
That will be a major goal of the protocols we will discuss

## Performance measures: Talking about the performance of a link

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- A common measure of link performance is the amount of time required to transfer a packet from one end to the other
- Delay = Propagation delay + Transmit delay + Queuing time
  - Propagation delay = time to send a signal end to end
  - Transmit delay = time to stream all bits of a packet onto the link
  - Queuing time = time the packet is delayed (buffer, queue, etc)

### An analogous situation: water in a pipe

- each gallon of water takes a finite amount of time to traverse the pipe, depending on the pipe's length (propagation delay)
- the pipe can deliver a certain number of gallons per second, depending on the pipe's size (transmit delay)
- water may be stored temporarily in a reservoir (queuing time)

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## Proagation delay (aka "latency")

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- = time it takes for the elictrical or optical signal to be sent from the sending end of the link to the receiving end
- Propagation delay =  $L / c$ 
  - $L$  = Length of the link (m)
  - $c$  = Speed of light in the medium the link is made from (m/s)
    - Free space:  $3.0 \times 10^8$  m/s
    - Copper:  $2.3 \times 10^8$  m/s
    - Optical fiber:  $2.0 \times 10^8$  m/s

### Example:

Latency of a 100Km copper link

$$\begin{aligned}
 &= 100 \times 10^3 \text{ m} / 2.3 \times 10^8 \text{ m/s} \\
 &= 10^{-3} / 2.3 \text{ s} \\
 &= 0.43 \times 10^{-3} \text{ s} = 430 \text{ usec}
 \end{aligned}$$

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## Transmit delay

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- = time it takes to stream the bits of the packet onto the link
- Transmit delay =  $p / bw$ 
  - $p$  = number of bits in the packet
  - $bw$  = bandwidth of the link (bps)

## Link bandwidth

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- Some types of link technologies and designs can handle more bits per second (bps) than others
- Bandwidth of a link = the specified max rate in bps at which the link can accept binary data
  - (Remember: rates expressed using powers of 10, not 2)

### Example:

If bandwidth = 100Mbps ( $10^2 \times 10^6$  bps =  $10^8$  bps) then:

- The link can accept  $10^8$  bits in a second, including both data and "overhead" bits
- When data is being streamed onto the link, a new bit can be handled each  $10^{-8}$  seconds (10 nsec)
- It will take  $n \times 10^{-8}$  sec to stream an  $n$ -bit packet onto the link (this is the Transmit delay for the packet)

## Queuing delay

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- = time packet is delayed in a buffer or queue
- Depends on the specific situation -- hard to generalize accurately
- We will normally ignore queuing delay in calculating link delay, but *don't forget it is there*.

## An example link delay calculation

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What is the minimum link delay for a 1MB packet sent from New York City to San Francisco (4800 Km) over a 10Mbps fiber link?

**Prop delay**       $= L / c$   
                      $= 4800 \text{ Km} / c_{\text{in\_fiber}}$   
                      $= 4.8 \times 10^6 \text{ m} / 2 \times 10^8 \text{ m/s}$   
                      $= \mathbf{24 \text{ msec}}$  (the "standard" US coast-to-coast delay)

**Transmit delay**    $= p / bw$   
                      $= (1 \times 2^{20} \text{ bytes} \times 2^3 \text{ bits/byte}) / (10 \times 10^6 \text{ bps})$   
                      $= 2^{23} / 10^7 \text{ sec}$   
                      $= \mathbf{839 \text{ msec}}$

**Delay**              $= 24 + 839 \text{ msec}$   
                      $= \mathbf{863 \text{ msec}}$

## Other performance metrics: RTT

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- Round-trip time (RTT) = network delay associated with sending a message and receiving a reply or acknowledgement
- $RTT = 2 \times \text{Propagation delay}$
- Note that this is a Best-Case time that considers the network delay only: it ignores processing and buffering times at each end
- Frequently used in rough performance measures

### Example

A Mars rover is controlled from Earth. What is the minimum lag time between the time that a rover sensor detects a problem and a responsive control command can be received?

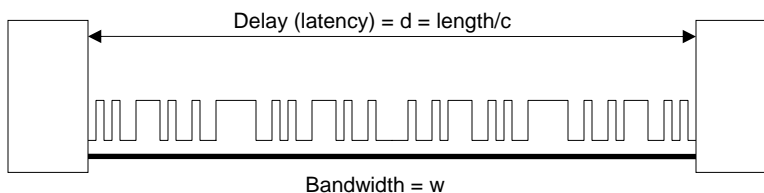
Assume Earth-Mars distance  $\approx 80 \times 10^6 \text{ Km} = 8 \times 10^{10} \text{ m}$   
 Then  $L = 8 \times 10^{10} / 3 \times 10^8 = 266 \text{ sec}$   
 $RTT = 532 \text{ sec} \approx 9 \text{ minutes}$

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## Another type of measure: How many bits can a wire “hold”?

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Sender starts a sending stream of bits at time  $t = 0$

The first bit reaches the receiver at time  $t = d$

The sender will send a bit each  $1/w$  seconds

So, by time  $t = d$ ,  
 the sender has sent  $d/(1/w) = dw$  bits

These  $dw$  bits are “in flight” somewhere between the sender and the receiver

### Example:

How much data stays in flight on a Coast-to-coast 100Mbps Fiber link?

Delay  $\approx 24\text{ms}$   
 $\# \text{ bits in flight} = 24\text{ms} \times 100\text{Mbps}$   
 $= 2.4 \times 10^6 \text{ bits}$   
 or about 290KB

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## delay x bandwidth product

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- = the maximum number of bits that can be in transit on the link simultaneously ("in flight")
- Sometimes used as a measure of the link's capacity
- Maintaining dxb bits in flight is called "keeping the pipe full"

### Example

A 100Mbps link with a latency of 10 msec carries streaming data from a sensor to a computer. When the computer's receive buffer is full, it sends a PAUSE signal to the sensor. Any data in flight at that time cannot be received and must be re-sent after the PAUSE is removed. What is the minimum amount of data that would have to be re-sent after a PAUSE?

$$\text{dxb product} = 100 \times 10^6 \text{ bps} \times 10 \times 10^{-3} \text{ sec} = 10^6 \text{ bits} = 122 \text{ KB}$$

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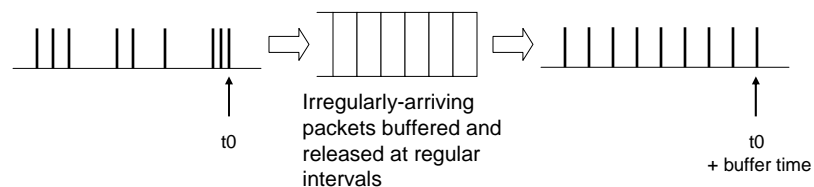
## Other performance metrics: jitter

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- Jitter = irregularity in delay between packets in a datastream



- A critical problem in streaming media (often resolved by buffering)



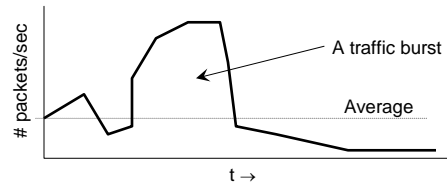
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## Traffic profiles

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- Most real network traffic is bursty:



- Knowing the average load is good, but is of limited usefulness -- we need to know the traffic "profile"
- Average loads usually don't kill networks, bursts kill networks